

Future Internet Transport Layer -Heading towards a Post-TCP Era?

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Outline

Status quo

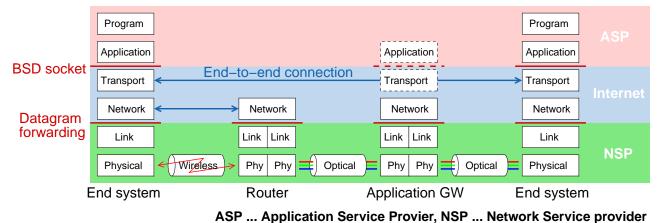
- Internet
- Internet transport layer
- TCP/IP midlife crisis
- Transport "layer" future
 - Evolution
 - Revolution
- Conclusions and outlook



Status Quo - Internet

Never standardized, still evolving...

Internet TCP/IP-"Architecture"



- Current Internet := layer 2.9 to layer 4.5
 - Network layer: Addressing, routing, peering
 - Transport layer: End-to-end interconnection of intelligent hosts
- Narrow APIs, no control-plane
- ➡ TCP/IP as "spanning" layer

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Status Quo - Internet

Internet Arpanet Design Paradigms

Design objectives

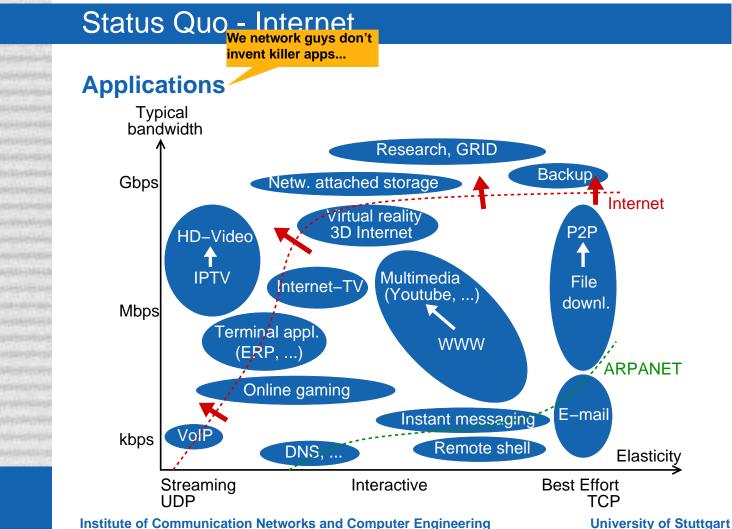
- 1. Communication must continue despite loss of networks or gateways
- 2. Multiple types of communications services
- 3. Accomodate a variety of networks
- 4. Distributed management of its resources
- 5. Cost effective
- 6. Host attachment with a low level of effort
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Paradigms

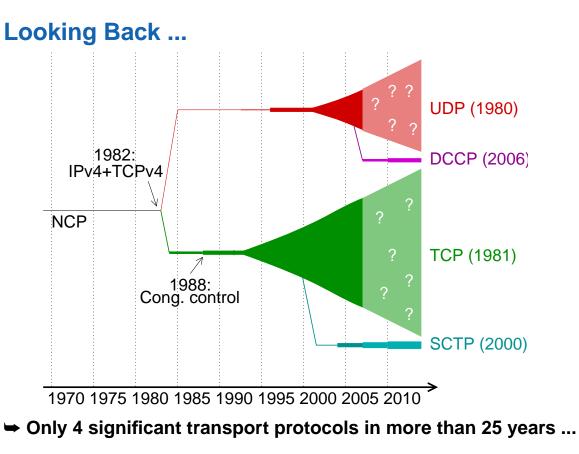
- 1. Packet switching
- 2. Layering (with simple APIs)
- 3. A network of collaborating networks (IP spanning layer)
- 4. Intelligence in end systems (end-to-end principle)

From D. Clark, "The Design Philosophy of the DARPA Internet Protocols", Proc. SIGCOMM 1988, Sept. 1988

- ➡ TCP/IP heavily influenced by these paradigms
- → New paradigms? Other protocol architecture!

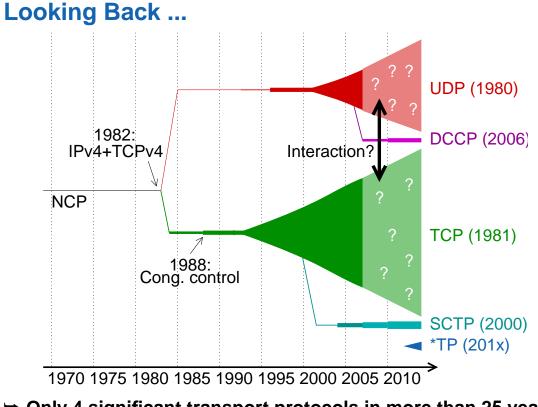


Status Quo - Internet Transport Layer



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Status Quo - Internet Transport Layer

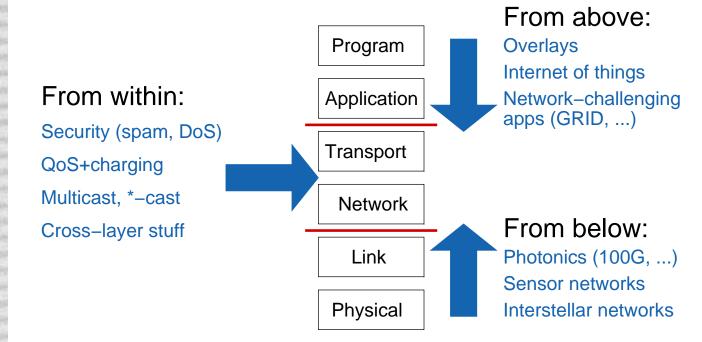


➡ Only 4 significant transport protocols in more than 25 years ...

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Midlife Crisis

New (and Old) Requirements



➡ "Future Internet: Just more QoS and multicast?" (H. Schulzrinne)

Midlife Crisis I: Congestion Control

Design Principles of V. Jacobson's Congestion Control

- Sender-side control of data rate by congestion window
- Greedy probing of available bandwidth on path (window increase)
- Implicit congestion feedback by packet loss (window decrease)

Characteristics

Just putting more bandwidth might not help ...

- Key impact on application performance
- Best effort, elastic applications only, no QoS
- Vague notion of fairness (unfair to connections with larger RTT)
- ➡ Never perfect, but usually good enough

However ...

- Network characteristics changed a lot since 1988
- More and more network-demanding applications
- Fairness may become an issue (e.g., network neutrality debate)

Midlife Crisis I: Congestion Control

Example: High-Speed WANs

Problem: Large bandwidth-delay products

- Standard TCP congestion control not well suited for large window sizes
 - Example: 10 Gbit/s TCP throughput with MTU=1500-byte, RTT=100 ms requires
 - Average congestion window of 83,333 segments
 - At most one drop/mark every 5 bill. packets (one drop every 1 2/3 hours)
- Long convergence times and significant unfairness
- Solution approaches
 - 1. Increase MTU
 - 2. UDP-based protocols (with some application-level congestion control)
 - 3. Modification of TCP congestion control in sender
 - More aggressive window increase functions
 - Delay-based congestion control
 - ➡ Multitude of proposals (Highspeed TCP, Scalable TCP, HTCP, FAST TCP, BIC TCP, CUBIC TCP, Compound TCP, ...)
- ➡ Next challenge: 100 Gbit/s links?

Midlife Crisis I: Congestion Control

Open Issues

• Extrem variety of networks

From sensor networks to high-speed optical networks

• Large range of application requirements

Many non-elastic applications (e.g., multimedia, pseudo-wires)

• TCP-unfriendly path characteristics

Variable link capacities, corruption packet loss, packet reordering, ...

Multi-domain operation

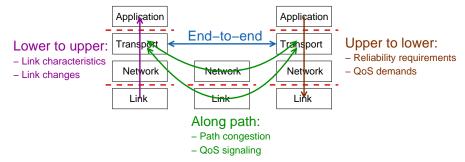
Misbehaving senders, receivers, and applications

- Fairness
- ...
- ► Any chance for *one* Internet congestion control?

Further details in: Michael Welzl, Dimitri Papadimitriou, Michael Scharf, "Open Research Issues in Internet Congestion Control", IETF internet draft, work in progress, July 2007, draft-irtf-iccrg-welzl-congestion-control-open-research-00.txt

Midlife Crisis II: Cross-Layer Issues

Cross-Layer Information Exchange



Significant potential for optimization

Challenges

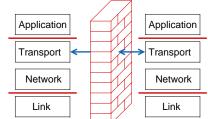
- Host local: Standardized interfaces (device OS application)
- Remote along path
 - Protocol extensions or new protocols (in-band or out-of-band)
 - Interaction with routing, IP tunnels, ...
 - Security issues (AAA, DoS prevention, ...)
- Internet control plane?

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Midlife Crisis III: End-to-end Paradigm

The Raise of Application Gateways

- Intermediaries break up end-to-end semantics
- Many functions
 - Security, information hiding (NAT/firewalls, session border controllers, ...)
 - Performance optimization (Web caches, WAN accelorators, ...)



Application GW

- Content processing (compression, virus checking, transcoding, ...)
- Facilitate rendezvous and/or forwarding (SIP proxies, SMTP relays, ...)

Problems

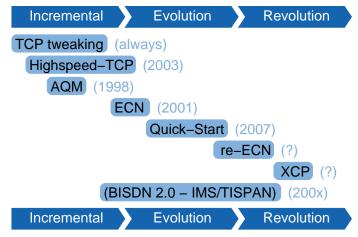
- May become single point of failure
- May hinder/limit communication and/or new protocol extensions
- May require certain trust relationships
- End-to-end vs. "balkanization" of the Internet?

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A bug, or a feature?

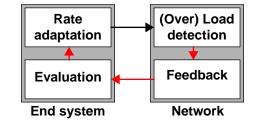
Transport "Layer" Future

Recent TCP Research+Standardization



- Major focus: Congestion control
- Main objectives
 - Optimize performance for high-speed and wireless networks
 - Increase fairness
- ➡ Question: Role of network, i. e., routers?

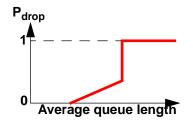
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Transport Layer Evolution: Examples

Active Queue Management (AQM) - RFC 2309 (1998)

- Replace drop-tail buffers in routers
 - Avoids synchronization effects
 - Improves fairness
- Reality check
 - Enabled in some routers
 - No single optimal parameter set



Simple idea, but still too complex for the real world?

Explicit Congestion Notification (ECN) - RFC 3168 (2001)

- Congested router sets bits in IP header instead of dropping packets
- Reality check
 - Support by major operating systems, but not enabled by default
 - Deployment problems with buggy routers and middleboxes
- Chicken-egg deployment problem and no pressing need?

Transport Layer Evolution: Examples

Quick-Start TCP Extension - RFC 4782 (2007)

Slow-Start:



- One pillar of TCP congestion control
- Exponential window growth

• Speeds up interactive WAN applications

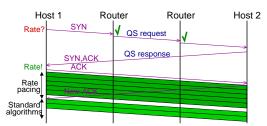
- After connection setup or idle periods
- For large bandwidth-delay products
- Reality check
 - Requires support in all routers
 - Some open (research) issues
- Short-term deployment in puplic Internet unrealistic

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Quick-Start:



- Recent experimental TCP extension
- (Almost) immediately use large window



Transport Layer Evolution: Examples

eXplicit Control Protocol (XCP) - RFC 5xxx

• Explicit congestion feedback from routers

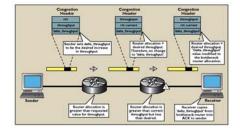
- Some congestion state in packets
- Feedback on rate increment/decrement
- Routers do some per-packet calculations (but no per-flow state)

Potential advantages

- High link utilization for high-speed WANs
- Fair bandwidth allocation

Reality check

- Complexity: Requires per-packet computations in routers
- Unsolved issues with short-lived flows
- Only a congestion control framework no transport protocol (so far)
- ➡ Revolutionary research ...



Source: A. Falk et. al., "Transport Protocols for High Performance"

Transport "Layer" Revolution

Internet Arpanet Design Paradigms

Design objectives

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But, any internetwork architecture has to provide

That's what transport layer is about ...

communication between various applications on end-systems

... not necessarily a "layer" with service primitives we have today

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Transport "Layer" Revolution

Actually, a network layer problem!

Design Space for Congestion Control

Implicit network feedback	or	Explicit network feedback
Loss-based, delay-based, bandwidth estimation techniques		In-band signaling Out-of-band signaling
No state in routers ("end-to-end")	or	Some state/processing in routers (per packet/per RTT/)
TCP friendly	or	More/less aggressive/fair/ than TCP
Best effort	or	Support of precedence/QoS

Design Space for Transport Functions

see SCTP and DCCP

Reliable transfer	or	Also partial reliable/unreliable transfer
Single stream	or	Multiple streams
.		•
Single path	or	Multiple (disjoint) paths
Unicast only	or	*cast support

(adapted from: S. Shalunov et. al., "Design Space for a Bulk Transport Tool")

Revolution - Food For Thought (1)

About Layering

Program	Connection-oriented (e.g., WS-*)
Application	Connection–less (HTTP)
Transport	Connection-oriented (TCP)
Network	Connection-less (IP)
Link	Connection–oriented (MPLS, ATM, UMTS,) Connection–less (Ethernet,)
Physical	Connection-oriented???

- Reduce "self-similarity" in stack?
- Get rid of static layering?
- Transport layer functions in user space, instead of kernel space?
 - Higher performance (locking, caching)
 - Less complexity, more flexibility
- What about better support for tunneling?

Revolution - Food For Thought (2)

About Granularity

- What is the minimum granularity of data exchange? What timescales?
- What are the communication primitives (in particular, for non-bulk-data transport)?
- Role of photonics (e.g., dynamic establishment transparent optical paths)?

Optical circuit	
Virtual path	
Connection	
Message	
Bit	

About Performance

- Where (and how) to handle resource sharing?
- Traffic engineering per aggregate/flow/connection/message/...? Centralized or distributed?
- What levels of reliability, availability, resilience?
- What about self-optimization?

Revolution - Food For Thought (3)

About Addressing

- Anything more intelligent than port numbers?
- Incorporate security and access control features?
- Handle heterogenity in naming and addressing?
- Connection-centric vs. data-centric?
- Support discovery and rendevous services?

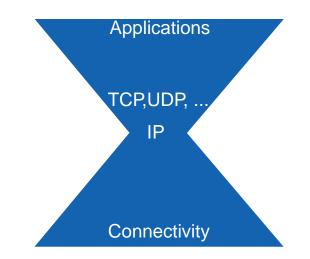
About Interfaces

- What degree of transparency?
- What interaction between data/control/management plane?
- What about peering interfaces? Role of business issues?

Conclusions and Outlook

Heading towards a Post-TCP Era?

- Transport layer functions are a key challenge for any "Future Internet"
- Evolution vs. revolution?
 - Short-term: TCP's shortcomings become more and more evident
 - Mid-term: Role and functions of routers to be rediscussed
 - Long-term: We will design Post-TCP once we know the "Future Internet"
- However ...

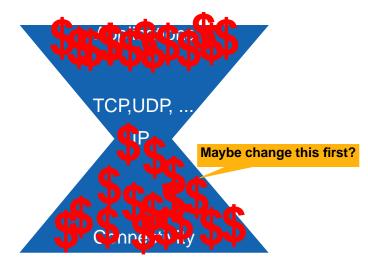


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