# **Trust Modeling**

### **Reasoning with Uncertainty**

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## Trust Modeling

#### **Outline**

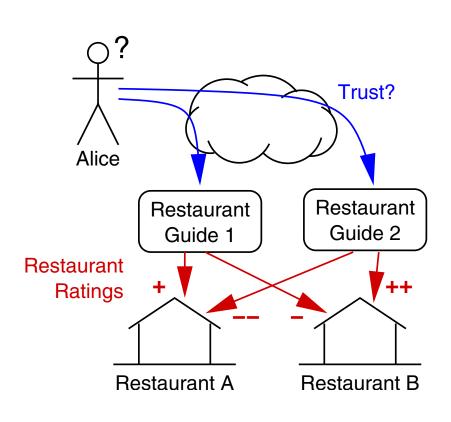
- Motivation
- Trust Modeling
  - Trust Relations
  - Reasoning with Trust
  - Representation of Trust Values
  - Trust Computation
- Proposal for a New Trust Model
- Conclusion and Outlook

### **Example: Restaurant Guides**

- Restaurant guide web services
- Problem
  - different restaurant guides may provide different results
  - anyone can offer a restaurant guide and disseminate falsified ratings
- ⇒ "Whom can I trust?"

#### **Trustworthiness**

- Competence ("is able to ...")
- Benevolence ("is willing to ...")
- **→** Need estimation of trustworthiness, e.g. for
  - decision whether or not to use a service
  - weighted combination of ratings



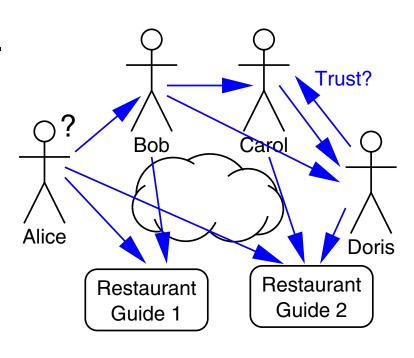
### Motivation

#### First-hand knowledge

- Good / bad own experiences, technical knowledge, guarantees, ...
- **→** But: often only for few services available!

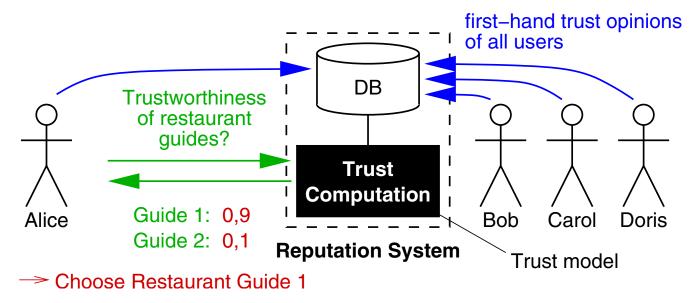
### Second-hand knowledge

- Exchange and evaluate trust estimations of other users
- → Again: "Whom can I trust?"
- Malicious / incompetent users
- Conflicting opinions, uncertainty, ...
- **▶** Need estimation of trustworthiness of trust estimations
- **→ Complex** graphs of trust relations, "Web of Trust"



### Goal

#### **Reputation System**



- All users publish (possibly false) first-hand trust opinions about other users and services
- Reputation system computes trustworthiness of any user / service

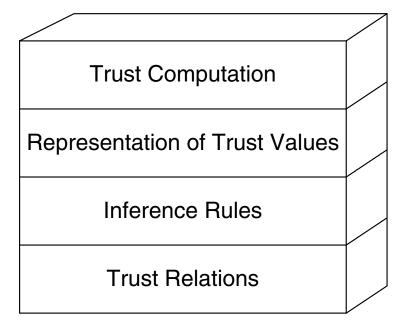
#### Note:

Reputation system do not aim to create or increase trust, nor to emulate (possibly irrational) human behaviour, but to serve a basis for a risk estimation.

## Trust Modeling

#### **Questions to answer**

- Nature of trust relations (properties)
- Reasoning with trust relations (inference rules)
- Representation of trust values (trustworthiness)
- Trust computation (trustworthiness of derived trust relations)



### Nature of Trust Relations

### **Working Definition**

- Trust is a unidirectional relation from truster to trustee, expressing the belief of the truster that the trustee will behave as expected.
- Distinguish between
  - direct (functional) trust: "Trustee has this property."
  - indirect (recommender) trust:
     "Trustee can recommend someone who has this property."
    - limit of recommendation hops

### **Trust Properties**

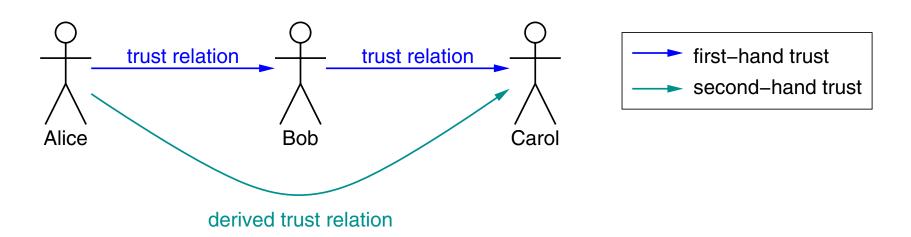
- Trust is specific to a given property / context
- Trust is not symmetric
- Trust is not reflexive
- Trust is not transitive in general
  - "A trusts B" and "B trusts C" does not necessarily imply "A trusts C"
  - must be specified in inference rules

## Reasoning with Trust

- Set of inference rules defining
  - Which trust relations can be derived from a set of existing trust relations?
- Example: Recommendation rule [A. Jøsang]

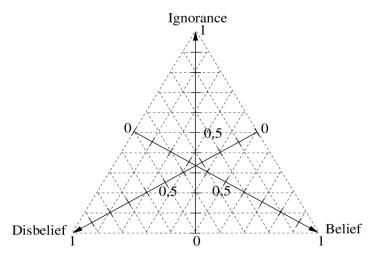
concatenation of two trust relations:

 $trust(Alice, Bob) \land trust(Bob, Carol) \Rightarrow trust(Alice, Carol)$ 



## Representation of Trust Values (Trust Metrics)

- Range: "distrust" ↔ "no trust" ↔ "trust"
  - in open systems: negative trust values often not useful
- Default value:
  - in open system: choose lowest possible value
- Uncertainty required?
- Granularity:
  - discrete values, e.g."no trust", "marginally trust", "full trust"
  - continuous, e.g.trust ∈ [0...1]
  - multi-value:trust ∈ [-1...1], confidence ∈ [0...1]
  - upper and lower bound / opinion triangle

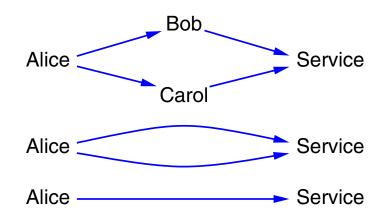


From: Audun Jøsang, "Artificial Reasoning with Subjective Logic"

## **Trust Computation**

### **Operator-based Trust Computation**

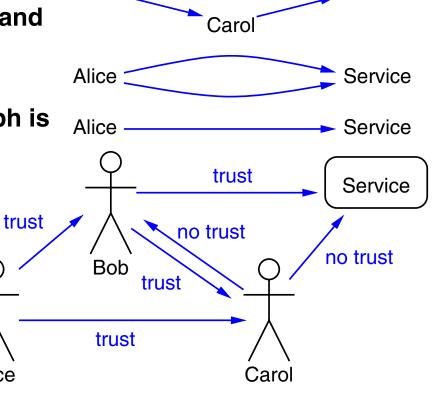
- Arithmetic operator for each combination rule
- Combining trust values of the involved trust relations
  - e.g. multiplication, min()/max(), average, fuzzy logic operators, ...
- Successive composition of serial and parallel trust relations



## **Trust Computation**

### **Operator-based Trust Computation**

- Arithmetic operator for each combination rule
- Combining trust values of the involved trust relations
  - e.g. multiplication, min()/max(), average, fuzzy logic operators, ...
- Successive composition of serial and parallel trust relations
- → Problem: only possible, if trust relation graph is a directed series-parallel graph



Bob

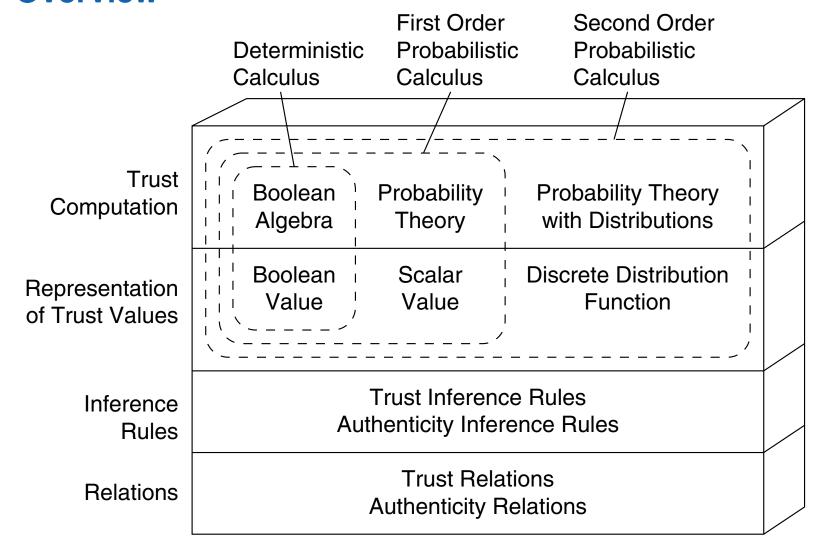
Alice

Alice

Service

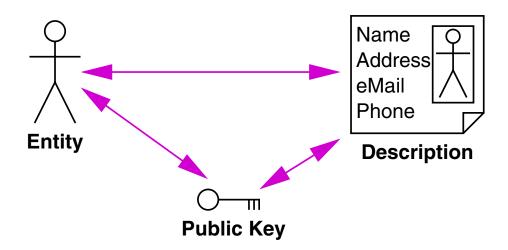
## Proposal for a New Trust Model

#### **Overview**



#### Why Authenticity Relations?

- Authenticity of exchanged trust opinions must be protected, e.g. with dititally signed trust certificates
- Recommendation systems used for authenticity validation of public keys (e.g., PGP Web of Trust)



#### **Trust Relations**

Relations (not signed)	Certificates (signed)
$ \begin{array}{c c} \hline  & c, h \\ \hline  & E_{A} \\ \hline  & E_{B} \\ \hline  & E_{B}, c, h) \end{array} $	
$ \begin{array}{c c} \hline  & c, h \\ \hline  & K_B \end{array} $ $ \begin{array}{c c} \hline  & K_B \end{array} $ $ \begin{array}{c c} \hline  & K_B \end{array} $ $ \begin{array}{c c} \hline  & E_A: Trust(K_B, c, h) \end{array} $	$ \begin{array}{cccc} & & & & & & & & \\ & & & & & & & \\ & & & & $
C, h  Address eMail Phone  D <sub>B</sub> E <sub>A</sub> :Trust(D <sub>B</sub> , c, h)	Name Address eMail Phone D <sub>B</sub> K <sub>A</sub> :Trust(D <sub>B</sub> , c, h)

E = entity

K = public key

D = description

c = context / property

h = recommendation hops

### **Authenticity Relations**

Relations (not signed)	Certificates (signed)
E <sub>A</sub> E <sub>B</sub>	
E <sub>A</sub> :Auth(K <sub>B</sub> , E <sub>B</sub> )	
E <sub>A</sub> Name Address eMail Phone  D <sub>B</sub>	
$E_A$ :Auth( $D_B$ , $E_B$ )	
K <sub>B</sub> — m  K <sub>B</sub> — m  Name Address eMail Phone  D <sub>B</sub>	Name Address eMail Phone
$E_A$ :Auth( $K_B$ , $D_B$ )	$K_A$ :Auth( $K_B$ , $D_B$ )

E = entity

K = public key

**D** = description

#### 12 Inference rules

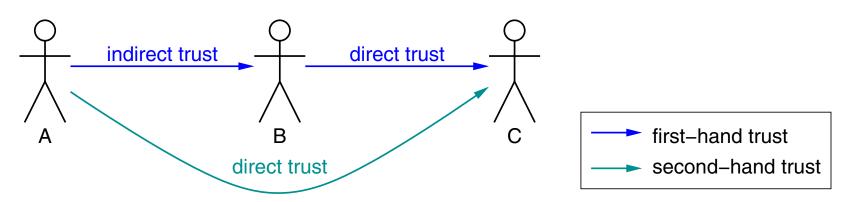
**Example 1: Transitive Trust Rule (2 parts):** 

1. indirect trust + direct trust ⇒ direct trust

A:Trust(B, c, h)  $\land$  B:Trust(C, c, 0)  $\land$  h>0  $\Rightarrow$  A:Trust(C, c, 0)

A, B: entity or public key

C: entity or public key or description



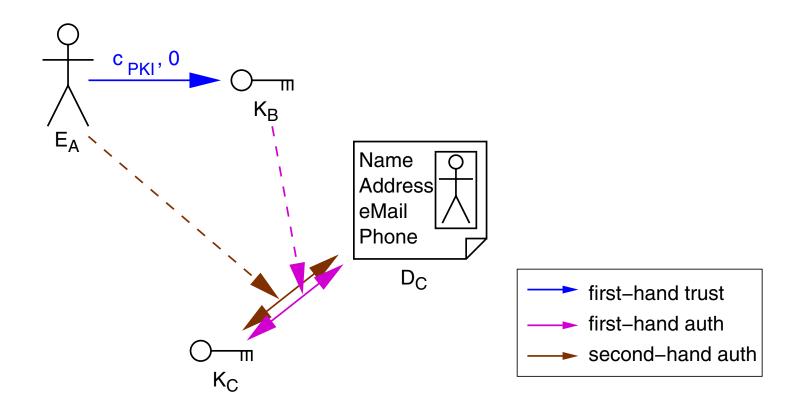
2. indirect trust + indirect trust  $\Rightarrow$  indirect trust

A:Trust(B, c,  $h_1$ )  $\land$  B:Trust(C, c,  $h_2$ )  $\land$   $h_1>1 <math>\land$   $h_2>0$   $\Rightarrow$  A:Trust(C, c, min( $h_1$ -1,  $h_2$ ))

#### **Example 2: Authenticity Inference with Identity Certificate Rule**

 $E_A$ :Trust( $K_B$ ,  $C_{PKI}$ , 0)  $\land K_B$ :Auth( $K_C$ ,  $D_C$ )  $\Rightarrow E_A$ :Auth( $K_C$ ,  $D_C$ )

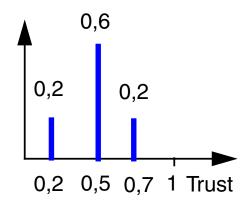
c<sub>PKI</sub>: property "issues valid identity certificates"



## Representation of Trust Values

### 3 Possibilities to represent trust values

- Boolean value: true / false very simple
- **2. Scalar Value:**  $t \in [0, 1]$  trust value interpreted as probability that the assumption is correct
- 3. Discrete distribution function
  allows to express uncertainty
  interpretation as second-order probability values



## **Trust Computation**

#### **Holistic Trust Computation**

- Interpretation of "trust" as "probability that the trustee has the named property"<sup>1</sup>
- **→** Trust values have well defined semantic
- **→** Computation with probability theory
  - → works for arbitrary trust structures!(in contrast to operator-based methods)

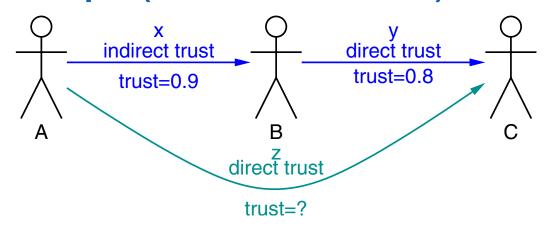
"Possible Worlds" Algorithm (for scalar trust values)

Each trust / authentication relation can be valid or invalid

- ⇒ 2<sup>n</sup> possible combinations (="possible worlds")
- 1. Check (for each "possible world"), whether the intended trust relation can be derived or not
- 2. Calculate the probability of occurence for each "successful" world
- 3. Resulting trust value = sum of probabilities of all "successful" worlds = probability of occurence of any "successful" world

<sup>1.</sup> Ueli Maurer, "Modelling a Public-Key Infrastructure"

#### **Example (scalar trust values)**



first-hand trust
second-hand trust

X	у	Z	probability
0	0	0	$(1-0.9)\cdot(1-0.8)$
0	1	0	$(1-0.9) \cdot 0.8$
1	0	0	0.9 · (1 – 0.8)
1	1	1	0.9 · 0.8

**Resulting trust value:**  $t = 0.9 \cdot 0.8$ 

(high computational complexity, more efficient computation algorithms exist)

### Conclusion and Outlook

#### Conclusion

- Reputation systems useful for various applications:
  - online auctions, PGP, P2P networks, ... (esp. for open user groups)
- Trust models must be designed carefully
  - distinguish direct and indirect trust
  - distinguish first-hand and second-hand trust estimations
  - be careful and precise with transitivity
- Operator-based trust computation → bad approach, better try holistic approach based on probability theory
- Integration of trust + authentication computation makes sense

#### Outlook

- Trust model evaluation
  - look out for counterintuitive effects → indicator for a bad model.
  - play attacker, try to fool your reputation system